



Revisiting SATZilla Features in 2024

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Abstract

Boolean satisfiability (SAT) is an \mathcal{NP} -complete problem with important applications, notably in hardware and software verification. Characterising a SAT instance by a set of features has shown great potential for various tasks, ranging from algorithm selection to benchmark generation. In this work, we revisit the widely used SATZilla features and introduce a new version of the tool used to compute them. In particular, we utilise a new preprocessor and SAT solvers, adjust the code to accommodate larger formulas, and determine better settings of the feature extraction time limits. We evaluate the extracted features on three downstream tasks: satisfiability prediction, running time prediction, and algorithm selection. We observe that our new tool is able to extract features from a broader range of instances than before. We show that the new version of the feature extractor produces features that achieve up to 26% lower RMSE for running time prediction, up to 3% higher accuracy for satisfiability prediction, and up to 15 times higher closed gap for algorithm selection on benchmarks from recent SAT competitions.

2012 ACM Subject Classification Theory of computation \rightarrow Logic and verification

Keywords and phrases Satisfiability, feature extraction, running time prediction, satisfiability prediction

Digital Object Identifier 10.4230/LIPIcs.SAT.2024.31

Supplementary Material *Software:* https://github.com/hadarshavit/revisiting_satzilla

Funding This research was partially supported by TAILOR, a project funded by EU Horizon 2020 research and innovation program under GA No. 952215. HH gratefully acknowledges support through an Alexander-von-Humboldt Professorship in Artificial Intelligence.

Acknowledgements We would like to thank Matthias König and Anja Jankovic for the discussions. We would also like to thank the reviewers who gave a valuable feedback.

1 Introduction

The Boolean satisfiability problem (SAT) is an important problem in computer science from both theoretical and practical viewpoints. Common usages of SAT include hardware and software verification, cryptography, and more. However, as SAT is \mathcal{NP} -complete, it takes substantial time and computing power to solve it, which becomes prohibitively expensive as formulas become larger. In order to deepen our understanding of SAT itself and develop better SAT solvers, it is of crucial importance to be able to describe SAT instances via an informative set of features.

Some of the most widely adopted such features are the SATZilla features [13, 7]. These are a fixed set of features that are calculated from the DIMACS CNF representation of a SAT instance. The SATZilla features consist of multiple groups, ranging from basic syntactic features describing the formula, such as its number of variables and clauses, to more complicated features, such as probing features derived from short runs of SAT solvers. Another type of features are based on statistics of graph representations of a given formula.



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27th International Conference on Theory and Applications of Satisfiability Testing (SAT 2024).

Editors: Supratik Chakraborty and Jie-Hong Roland Jiang; Article No. 31; pp. 31:1–31:26

Leibniz International Proceedings in Informatics



LIPICs Schloss Dagstuhl – Leibniz-Zentrum für Informatik, Dagstuhl Publishing, Germany

45 The SATZilla features have been successfully used in various domains, such as empirical
46 performance models (EPMs; also known as performance or running time prediction) and
47 algorithm selection [16, 17], algorithm configuration [6], and benchmark generation [8]. The
48 features are also used for caching in CDCL-based model counting solvers [14] and in a variety
49 of SAT solvers that incorporate machine learning techniques [4]. However, the SATZilla
50 features in their latest version date back to 2012. Since then, the SAT community has
51 undergone various changes. Most notably, SAT instances that we typically encounter today
52 have a larger number of variables and clauses, thus taking significantly more time and memory
53 to preprocess. Currently, the existing SATZilla feature extraction tool is unable to compute
54 many of these features because of time and memory limitations.

55 In this paper, we revisit the SATZilla features and introduce a new version of the feature
56 extraction tool. First, we replace the underlying solvers and the preprocessor with their
57 most up-to-date versions. We then fix compilation errors and other memory errors related
58 to dealing with larger formulas. Finally, we allow the user to set the time limits for feature
59 computation. We compare the performance of our new tool with the old one in two SAT
60 competitions. We measure the running times and the number of extracted features to check
61 for performance gains of our new tool. We then evaluate the extracted features on three
62 downstream tasks: satisfiability prediction, performance prediction and algorithm selection.
63 We show that our new tool yields an important advantage in performance compared to the
64 old tool across all three tasks.

65 The rest of the paper is organised as follows. We give a historical overview of the
66 development and applications of the SATZilla features in Section 2. We then introduce the
67 technical definitions, as well as the standard methodological pipeline in Section 3, wrapping
68 it up with the contributions of this study. The results are presented in Section 4, and
69 conclusions drawn in Section 5.

70 **2 Related work**

71 The SATZilla features were first introduced by Nudelman et al. [13] to construct EPMs, *i.e.*,
72 machine learning models that predict the running time of various SAT solvers given the
73 features representing SAT instances. The authors identified key features that contribute
74 the most towards having a good EPM prediction. Consequently, they used the EPMs as a
75 basis for algorithm selection, in which the algorithm that is selected corresponds to the one
76 with the lowest predicted running time. They leveraged the *performance complementarity*
77 phenomenon of SAT solvers, where no SAT solver dominates all others over all instances.
78 Therefore, selecting the best solver for each instance results in substantially better performance
79 compared to choosing any standalone solver for all instances. The SATZilla features were
80 also successfully used for the satisfiability prediction task for SAT instances from various
81 distributions [2].

82 Further developments in algorithm selection led to the 2007 version of the SATZilla
83 algorithm selector [16], which combined running time and satisfiability prediction with
84 other improvements. It won multiple medals in the SAT competition. This shows that
85 extracting features from SAT instances can also speed up the solving process, and not only
86 help understanding SAT. A newer version of the SATZilla algorithm selector was introduced
87 in 2012 [17], winning multiple awards as well. It used a random forest as a predictor, and
88 introduced an ensemble of pairwise classifiers to establish a ranking of the solvers, instead
89 of predicting the running times directly. Additional features were introduced thereafter,
90 revealing further important information about SAT instances.

Another common usage of the SATZilla features is (model-based) algorithm configuration. To this end, the hyperparameters of SAT solvers are optimised such that their performance is as good as possible on all instances. As running SAT solvers is computationally expensive, a surrogate model is employed; it takes as an input a configuration of hyperparameters and instance features and predicts the performance of the SAT solver using a given configuration on a given instance. The instance features boost the accuracy of the surrogate model. An example of that is SMAC [6], which successfully used the SATZilla features to optimise the performance of a wide range of SAT solvers on various benchmarks.

Last but not least, feature-based EPMs are proven to be useful for benchmark generation [8]. In this context, given a new instance, we compute the features, which is usually cheaper than running a SAT solver, and use the EPM to predict whether the instance is hard (or not) for the solvers at hand.

3 SATZilla features

The SATZilla features describe the SAT formula using various representations and statistics. We briefly introduce three graph representations of a SAT formula, as undirected graphs are a meaningful representation of SAT, maintaining the permutation invariance: a) variable graph: nodes are variables, an edge exists if variables appear in the same clause; b) clause graph: nodes are clauses, an edge exists when two clauses share a negated literal; and c) variable-clause graph: nodes are variables and clauses, an edge exists between a variable node and a clause node if the variable appears in the clause.

Feature computation starts with the *preprocessing* of the formula. This step, performed before solving an instance, renders the formula more accessible for SAT solvers. This means that the features are also computed on the version of the formula that is close to the one seen by the solvers. We believe that all these aspects can be boosted by using a modern preprocessor. Features are classically computed using the SATELITE preprocessor. We instead use the SBVA preprocessor, suggested by the winning solver from the 2023 SAT Competition. SBVA is also able to terminate after a set cutoff time, allowing for partial preprocessing, while SATELITE does not include this functionality. The preprocessed formula can then be directly used by a SAT solver without additional preprocessing within the solver, which improves the performance of algorithm selection.

Following the preprocessing, the *feature extraction* begins. There are ten feature groups that can be extracted. We note that we describe the feature groups according to their implementation in the SATZilla feature extraction tool, not according to their definitions from the corresponding paper [7]. We point out that all feature groups include the time required to compute the features in the group.

Preliminary features include the number of variables and clauses before/after preprocessing. The running time of this feature group includes the preprocessing time and the time required to read the formula. This group contains 7 features.

Basic features are cheap features that provide a basic description of the formula. They consist of the variable-clause ratio, the ratio between positive and negative literals in each clause, the number of unary, binary and ternary clauses, as well as statistics on clause nodes in the variable-clause graph. This group contains 15 features.

KLB are expensive features that include the node degree statistics of the variable nodes in the variable-clause graph, and the ratio of positive to negative occurrences of each variable. They also include measures for the proximity to Horn formula, such as the fraction of Horn clauses and statistics on the number of times each variable appears in a Horn clause. This

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137 group contains 21 features.

138 **Clause graph** (CG) features are expensive features that contain statistics on the degree of
139 the nodes in the clause graph, as well as the clustering coefficient. This group contains 11
140 features.

141 **Diameter** features contain information on the diameter of the variable graph, which is the
142 shortest path between each pair of nodes in the graph. This group contains 6 features.

143 **DPLL probing** (or unit propagation) features are computed by running the DPLL algorithm
144 for various depths and measuring the number of unit propagations at each depth. This group
145 contains 6 features.

146 **Lobjois** features are an estimation of the size of the search space. They are computed by
147 running the DPLL algorithm multiple times until a contradiction is found. Then, the average
148 depth of the contradictions is the log-estimation of the search space. This group contains 3
149 features, which are all based on the work of Lobjois and Lemaître [11].

150 **Survey propagation** features are based on computing statistics on the following probabilities
151 returned by the VARSAT [5] solver: a probability of each variable to be assigned to True, to
152 False, and to be unconstrained. This group contains 19 features.

153 **Clause learning** (CL) features are based on running a CDCL solver (ZChaff rand [12] in
154 the 2012 version, CadiCal [1] in our new version) for two seconds. We measure the number
155 and length of learned clauses for every 1000 decisions, and compute the statistics of those
156 values. This group contains 19 features.

157 **Local search** (LS) features are obtained by running two local search solvers many times,
158 each time up to 10000 steps, and computing statistics on those runs. In the 2012 version,
159 the local search solvers are GSAT and SAPS. We instead use GSAT and Sparrow 2011 in
160 our new version. This group contains 24 features.

161 **Linear programming** (LP) features are based on solving a relaxed version of the SAT
162 formula, where $C_1 \wedge C_2 \wedge \dots \wedge C_N$ is a Boolean formula with N CNF clauses $C_1 \dots C_N$
163 over Boolean variables x_j . We now consider linear programming variables x_j and solve the
164 following linear programming problem: Maximise $\sum_{i=1}^N \sum_{l \in C_i} v(l)$, where the value $v(l)$ of
165 literal l is defined as $v(x_j) = x_j$, $v(\neg x_j) = 1 - x_j$, while keeping $\sum_{l \in C_i} v(l) \geq 0$ for each C_i
166 and $0 \leq x_j \leq 1$ for each x_j . This means that every variable has a value between 0 and 1, each
167 clause has a value which is the sum of values of all literals, and the value of the formula is the
168 sum of values of all clauses. The goal is to maximise the value of the formula, while keeping
169 the value of every clause positive. Finally, statistics on the linear programming solution are
170 extracted. In the new version, we upgrade the linear programming solver package, lp solve.
171 This group contains 7 features.

172 We note that the basic, KLB and clause graph features are computed sequentially, with
173 each feature group being dependent on the successful computation of the previous groups;
174 *e.g.*, if the computation of the basic features fails, the KLB and clause graph features will
175 not be computed at all.

176 Another change we applied to the feature extraction tool is allowing a more precise
177 timeout setting. In the 2012 version, the time limits were hard-coded and set to high values
178 (for example, 1200 seconds for preprocessing). This can cause many feature computations to
179 simply terminate without computing any features at all. To this end, we adjust the code in
180 the new version to allow the user to set the time limits through a command line argument.

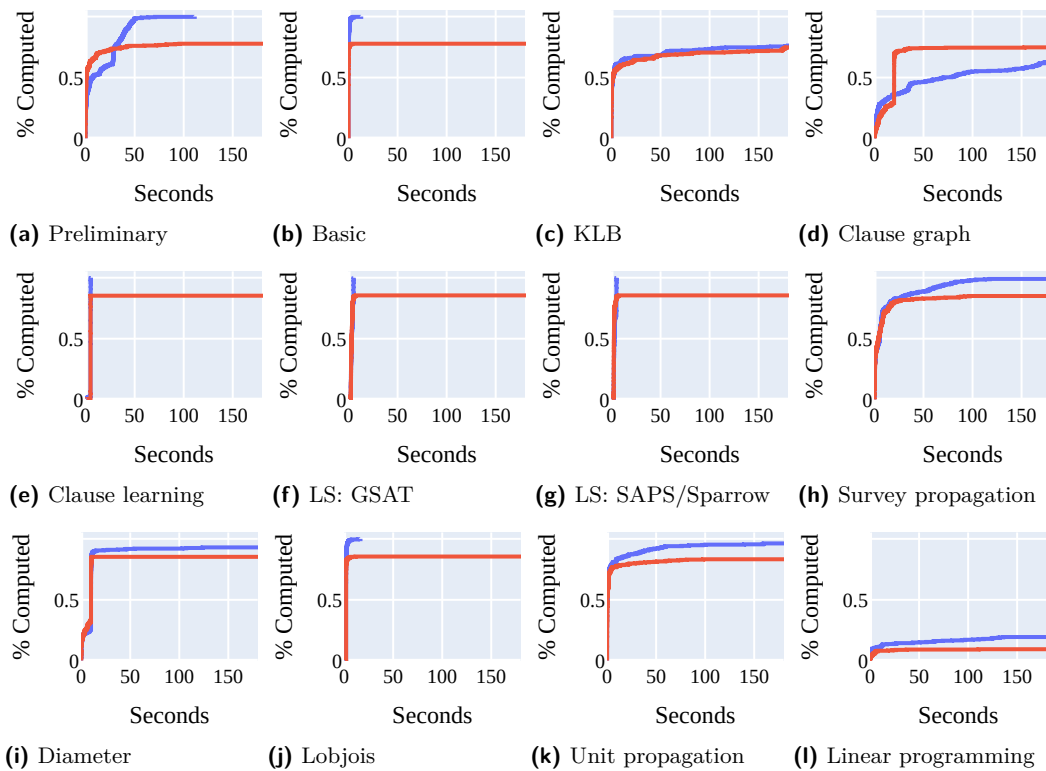
181 **4 Experiments**

182 We evaluate our new SATZilla feature extraction tool on the formulas stemming from two
 183 latest (2022 and 2023) SAT Competitions. We first look at the feature extraction times and
 184 then evaluate the features on three downstream tasks: satisfiability prediction, running time
 185 prediction, and algorithm selection. To assess the advantage of using the new version of the
 186 extraction tool, we extract the features using both our new version and the 2012 version of
 187 the tool, with a time limit of 180 seconds per feature group.

188 We use a cluster of 18 nodes, each equipped with 2 AMD EPYC 7543 32-core CPUs with
 189 256 MB L3 cache. Each node also has 1TB of memory. The cluster is running on a Rocky
 190 Linux 9.4 operating system. We measure running times using the runsolver tool [15].

191 **4.1 Feature computation time**

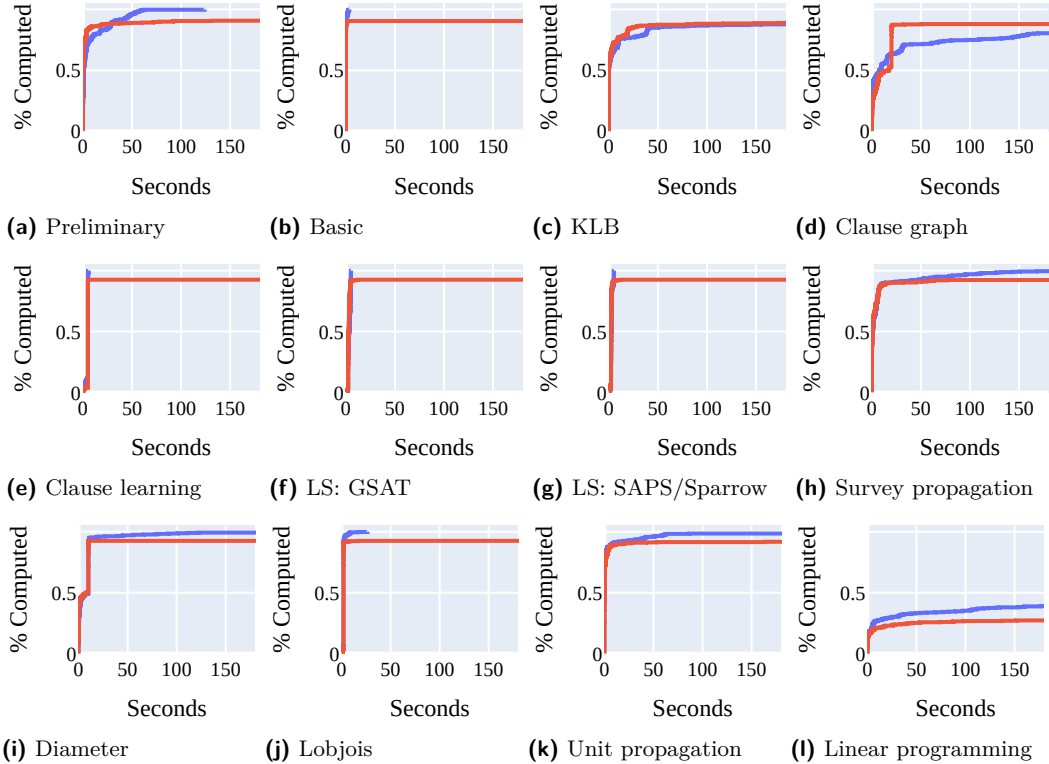
Figure 1 Percentage of features computed by the old tool (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue) over the available time budget for each feature group on the 2022 SAT Competition. For most feature groups, the new tool extracts features from more instances than the old one.



192 Figure 1 and Figure 2 show the percentage of features computed over the available time
 193 budget using the old (depicted in red) and new (depicted in blue) SATZilla tool on the 2022
 194 and 2023 SAT Competition data, respectively. For simplicity, due to the similarity of the
 195 overall results between the two competitions, we draw more detailed insights based on plots
 196 from the 2023 edition (Figure 2).

197 We first observe that the new tool is able to extract more features than the old one for

■ **Figure 2** Percentage of features computed by the old tool (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue) over the available time budget for each feature group on the 2023 SAT Competition. For most feature groups, the new tool extracts features from more instances than the old one.



198 most feature groups. In particular, we highlight the performance gains on the preliminary
 199 feature group (Figure 2a), for which the new tool can extract the features for all formulas,
 200 compared to less than 80% of the formulas when using the the old tool. We note that for
 201 some feature groups, like graph learning (Figure 2d), the old tool is able to extract more
 202 features compared to the preliminary feature group. This is due to the fact that the old tool
 203 extracts the preliminary, basic, KLB and CG feature groups together. Therefore, in case
 204 computing one of those groups takes a long time, the whole feature extraction fails.

205 We also note that for KLB and CG features there is an advantage for the old version,
 206 due to the new SBVA preprocessing method yielding larger formulas than its predecessor
 207 SATELITE (as in many cases smaller formulas are not always easier to solve). Similarly, the
 208 expensive graph-based features (*e.g.*, Figure 2d and Figure 2i) require more time to extract
 209 than smaller formulas. This is more apparent in the 2022 SAT Competition, where larger
 210 instances were used than in the 2023 SAT Competition.

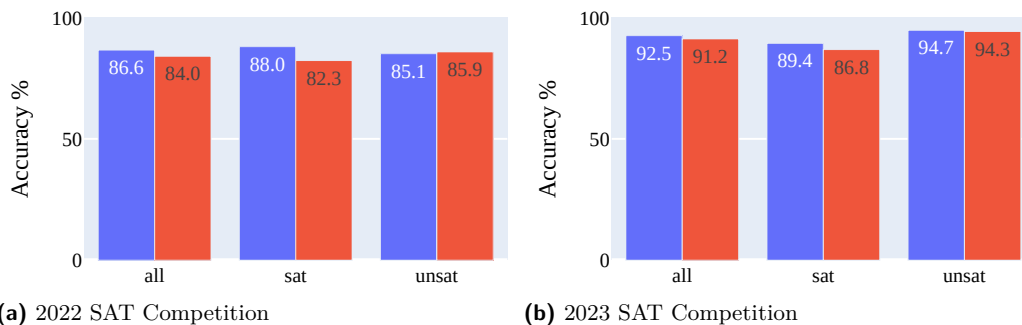
211 4.2 Satisfiability prediction

212 The first downstream task is satisfiability prediction, for which we measure the performance
 213 when using features extracted by our tool. We use the random forest (RF) classifier from the
 214 scikit-learn package to learn the mapping between features (representing SAT instances) and
 215 outputs (satisfiable or unsatisfiable). We optimise the hyperparameters of the RF for one

216 hour using SMAC3 [9] on 10-fold cross-validation of the training data. We consider instances
 217 from the 2022 and 2023 SAT Competitions for which we know the satisfiability result (put
 218 differently, we omit instances for which the solution is unknown). On each competition, we
 219 evaluate the performance of the RF model using 10-fold cross-validation. This results in
 220 having outer cross-validation (for evaluation) and inner cross-validation (for training). Such
 221 techniques have been previously used by AutoFolio [10].

222 We present the satisfiability prediction accuracy scores in Figure 3. We see that, by
 223 using features extracted via the new tool, we achieve better performance across all instances
 224 on both SAT competitions. Furthermore, we notice that the new tool leads to a higher
 225 accuracy gain for satisfiable instances than for unsatisfiable instances. For the latter, the
 226 accuracy remains very similar to the one achieved by using the old tool in the 2023 SAT
 227 Competition and slightly dropped for the 2022 SAT Competition. This might be due to
 228 the fact that the unsatisfiable instances are larger on average, thus being more prone to
 229 timeouts even when using our new tool, which goes along with the worse performance in
 230 the 2022 SAT Competition, where the unsatisfiable instances were larger than in the 2023
 231 SAT Competition. We point out that such high accuracy was already achieved before for
 232 industrial SAT instances [2].

■ **Figure 3** Accuracy of the satisfiability prediction task using a random forest with features extracted by the old (SATZilla 2012) and the new tool (SATZilla 2024). We see an overall higher accuracy for the new tool, which results from higher accuracy on satisfiable instances. On unsatisfiable ones, the accuracy remains the same.



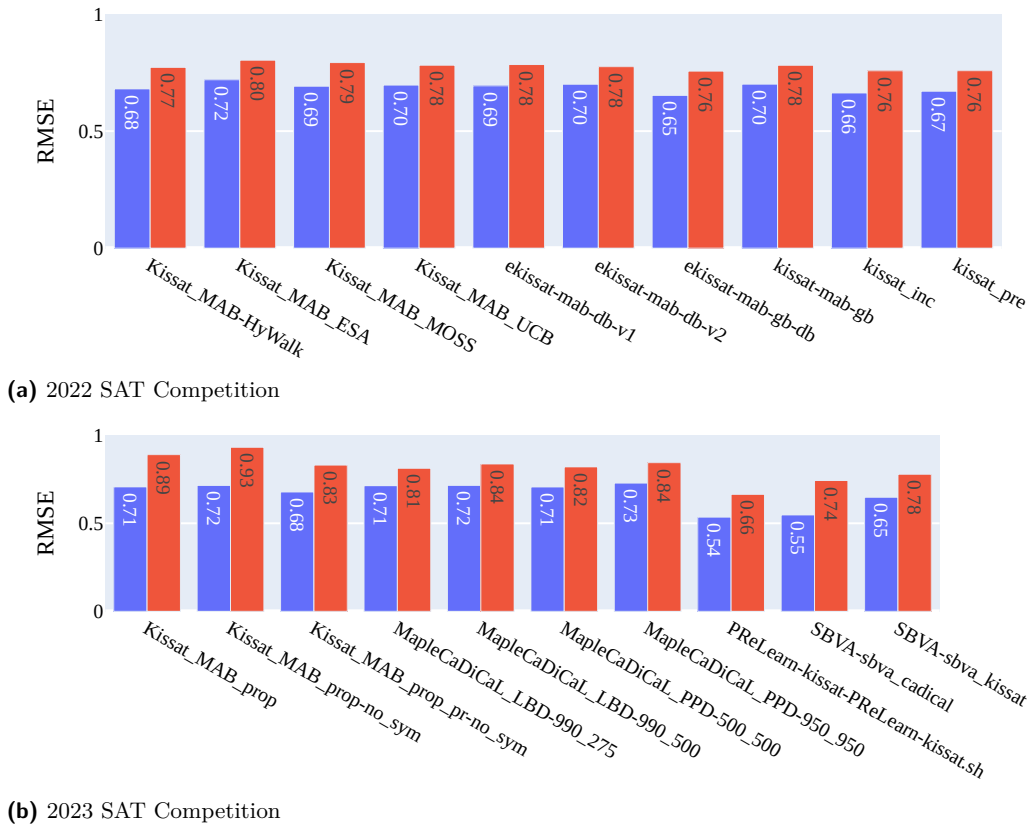
233 4.3 Performance prediction

234 The second downstream task we investigate is performance prediction, which has important
 235 applications in algorithm selection, configuration and benchmark generation. We refer to the
 236 methodology described in [7] and use a RF regressor as the EPM. It is important to note that
 237 for an accurate running time prediction we need to perform a \log_{10} transformation of running
 238 times prior to training the model, as done by Hutter *et al.* [7]. This transformation allows to
 239 capture the order of magnitude rather than small variations in the running time. The EPM
 240 then maps instance features to the log-transformed running times, and we aim to minimise
 241 the root mean squared error (RMSE) of the model, which is defined as $\sqrt{\frac{1}{n} \cdot \sum_{i=1}^n (\hat{y}_i - y_i)^2}$
 242 for n predicted running times \hat{y}_i and ground truth running times y_i . Lower RMSE scores are
 243 better and 0 is the optimal value. In line with the previous task, we perform inner and outer
 244 cross-validation and optimise the RF's hyperparameters for one hour using SMAC3.

245 We look into the performance of the EPM for running time prediction on all solvers from
 246 the 2022 and 2023 SAT Competitions. Here, we do not actually run solvers on instances to

247 record their running times, but rather use the running times reported by the competitions.
 248 We display the results for the 10 best solvers from each competition in Figure 4 (the results
 249 for all solvers can be found in the supplementary material). We see that using the features
 250 extracted by the new tool leads to the lower RMSE for all solvers, compared to using those
 251 extracted by the old tool. For some solvers, we observe a significantly lower RMSE, like
 252 `Kissat_MAB_prop-no_sym`, where using the new tool decreases the RMSE from 0.84 to
 253 0.69. Figure 5 shows the histogram of the error percentage for all solvers in the 2022 and 2023
 254 SAT Competitions. We see that for the 2022 competition, the new tool has more instances
 255 with error rate lower than 10% compared to the old version. For the 2023 competition, using
 256 the features extracted with the new version, more instances are predicted with less than 1%
 257 error. Histograms per solver are available in Appendix B.

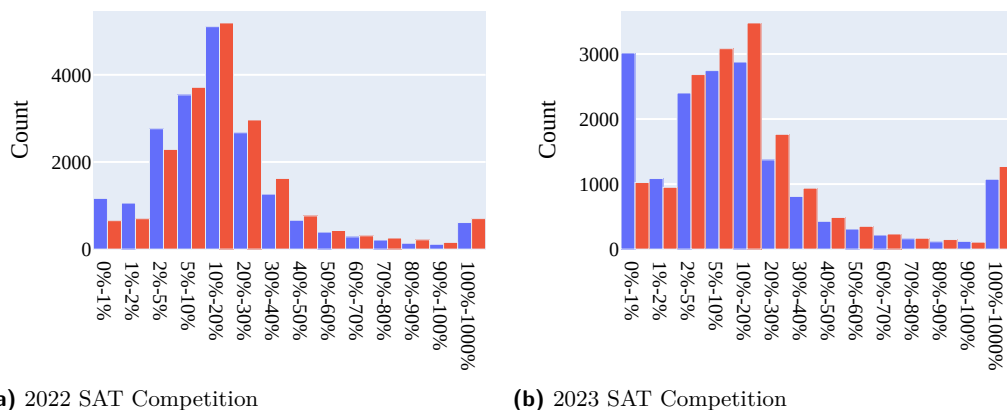
■ **Figure 4** Root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. The new tool achieves lower RMSE for all solvers.



258 4.4 Algorithm selection

259 The third and final downstream task we consider is algorithm selection. In algorithm selection,
 260 given a set of instances I , a set of solvers (*i.e.*, algorithm portfolio) \mathbf{A} and a performance
 261 metric $m : \mathbf{A} \times I \rightarrow \mathbb{R}$, we build an algorithm selector $S : I \rightarrow \mathbf{A}$ such that its performance
 262 is optimal on the instance set I according to the metric m . We compare the performance
 263 of algorithm selection to two standard baselines, the single best solver (SBS; *i.e.*, the solver

■ **Figure 5** Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. The new tool achieves lower error percentages.



with the lowest overall running time) and the virtual best solver (VBS; *i.e.*, the theoretical oracle which for each instance selects the actual best solver on it).

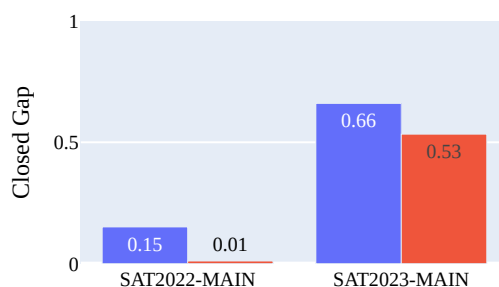
We measure the performance of algorithm selection using *closed gap*, which is computed as $\frac{m_{SBS} - m_S}{m_{SBS} - m_{VBS}}$, (*i.e.*, the closed gap stands for how much of the gap between the SBS and the VBS is closed by using the algorithm selector). We then use AutoFolio [10] as an algorithm selector, which consists of multiple algorithm selection approaches, from which the best one is suggested by using algorithm configuration. As algorithm portfolio for the selection, we use the 10 best solvers from each SAT competition. We train the selector for 8 hours.

Figure 6 shows the closed gap results on the 2022 and 2023 SAT Competitions. Positive closed gap values on both scenarios using both the old and the new tool indicate that, in general, SATZilla features are useful for the algorithm selection task. Importantly, features extracted with the new tool lead to better closed gap values on both scenarios. In Figure 7, we provide ECDF plots showing the fraction of instances solved over time. In the 2022 SAT Competition scenario, the old version of the tool performs worse than the SBS until a budget of approximately 1000 seconds, while the new version of the tool performs similarly to the SBS. After 1000 seconds, both versions of the SATZilla features perform similarly. In the 2023 SAT Competition scenario, the new tool performs better than the old one for budgets between 100 and 1000 seconds. With a budget of less than 10 seconds, the old tool solves more instances. Overall, the ECDF plots reflect well what is shown in Figure 6, where we see that the new tool exhibits a few percents higher closed gap value.

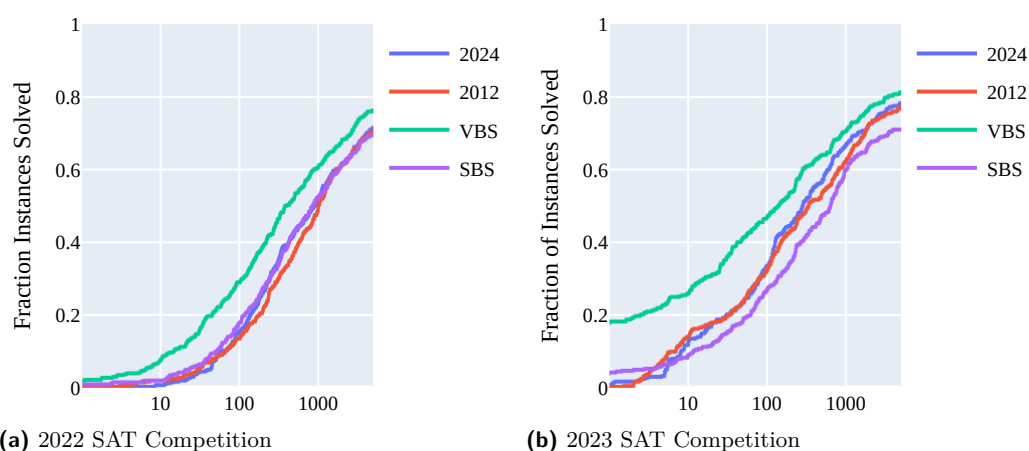
5 Conclusions

In this paper, we introduced an improved version of the well-known SATZilla feature extraction tool, motivated by the need to facilitate the feature extraction process by incorporating better user infrastructure and bug fixing. Our new version uses most up-to-date preprocessing techniques and SAT solvers, which allow for better representation of SAT formulas. Our experiments showed that, by using the new tool, we achieve a more accurate satisfiability prediction, a lower error for running time prediction, and a better closed gap for algorithm selection.

■ **Figure 6** Closed gap values for the algorithm selection task using the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue) on the 2022 and 2023 SAT Competitions; higher is better.



■ **Figure 7** ECDF plots for the algorithm selection task using the old (SATZilla 2012) and the new tool (SATZilla 2024) on the 2022 and 2023 SAT Competitions.



292 Our new SATZilla 2024 extraction tool aims to facilitate and promote the usage of SAT
 293 features even beyond their current scope. It is easily extensible and thus encourages building
 294 atop, *e.g.*, by looking into features based on the recent developments in the explainability of
 295 SAT solvers [3], or other advancements in SAT.

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356 **A** Running time prediction results

357 In this appendix, we present the full results of running time prediction for all solvers from
358 the 2022 and 2023 SAT Competitions. In Table 1 and Table 2 we show the results for the
359 2022 SAT Competition, where we see that our new tool constantly outperforms the old one.

■ **Table 1** RMSE of random forest for predicting log-transformed running times of SAT solvers from the 2022 SAT Competition using the old and new SATZilla features.

Solver	SATZilla 2012	SATZilla 2024
CaDiCaL-watchesat-lto	0.70	0.63
CaDiCaL_DVDL_V1	0.73	0.65
CaDiCaL_DVDL_V2	0.75	0.65
CadicalReorder	0.75	0.66
Cadical_ESA	0.75	0.66
IsaSAT	0.73	0.64
Kissat-MAB-rephasing	0.74	0.67
Kissat_MAB-HyWalk	0.77	0.68
Kissat_MAB_ESA	0.80	0.72
Kissat_MAB_MOSS	0.79	0.69
Kissat_MAB_UCB	0.78	0.70
Kissat_adaptive_restart	0.72	0.68
Kissat_cfexp	0.79	0.71
LSTech_CaDiCaL	0.77	0.63
LSTech_Maple	0.74	0.66
LSTech_kissat	0.76	0.67
LSTech-Maple-BandSAT	0.67	0.62
LSTech-Maple-FPS	0.74	0.65
LSTech-Maple-HyWalk	0.72	0.64
MapleLCMDistChrBt-DL-v3	0.61	0.58
MergeSat 4.0-rc-rc3	0.65	0.61
SLIME SC-2022	0.70	0.65
SLIME SC-2022-alpha	0.70	0.66
SLIME SC-2022-beta	0.71	0.66
SLIME SC-2022-gamma	0.73	0.69
SeqFROST-ERE-All	0.76	0.68
SeqFROST-NoExtend	0.77	0.67
cadical-hack-gb	0.73	0.62
cadical_rel_Scavel	0.71	0.61

360 Similarly, in Table 3 we show the results for the 2023 SAT Competition, where using the
 361 features extracted by our new tool leads to lower RMSE than with the old one.

362 **B** Running time prediction histograms

363 We present the histograms for the RMSE values per solver for running time prediction in
 364 Figures 8-19.

■ **Table 2** RMSE of random forest for predicting log-transformed running times of SAT solvers from the 2022 SAT Competition using the old and new SATZilla features (contd.).

Solver	SATZilla 2012	SATZilla 2024
ekissat-mab-db-v1	0.78	0.69
ekissat-mab-db-v2	0.78	0.70
ekissat-mab-gb-db	0.76	0.65
glucose-reboot	0.78	0.68
hCaD_V1-psids	0.69	0.63
hCaD_V2	0.68	0.64
hKis-psids	0.69	0.59
hKis-sat	0.79	0.69
hKis-unsat	0.74	0.66
kissat-els-v1	0.76	0.67
kissat-els-v2	0.74	0.67
kissat-els-v3	0.78	0.67
kissat-els-v4	0.77	0.66
kissat-mab-gb	0.78	0.70
kissat-sc2022-bulky	0.76	0.69
kissat-sc2022-hyper	0.76	0.70
kissat-sc2022-light	0.75	0.69
kissat-watchesat-lto	0.73	0.67
kissat_inc	0.76	0.66
kissat_pre	0.76	0.67
kissat_relaxed	0.72	0.67

■ **Table 3** RMSE of random forest for predicting log-transformed running times of SAT solvers from the 2023 SAT Competition using the old and new SATZilla features.

Solver	SATZilla 2012	SATZilla 2024
AMSAT_	0.80	0.77
BreakID-kissat-low.sh	0.92	0.83
CaDiCaL_vivinst	0.84	0.73
Cadical_ESA	0.85	0.74
Cadical_rel_1.5.3.Scavel	0.80	0.74
IsaSAT	0.90	0.82
Kissat_Inc_ESA	0.88	0.73
Kissat_MAB_Binary	0.89	0.72
Kissat_MAB_Conflict	0.87	0.72
Kissat_MAB_Conflict+	0.88	0.75
Kissat_MAB_DeepWalk+	0.87	0.73
Kissat_MAB_ESA	0.91	0.74
Kissat_MAB_Rephases	0.84	0.71
Kissat_MAB_prop	0.89	0.71
Kissat_MAB_prop-no_sym	0.93	0.72
Kissat_MAB_prop_pr-no_sym	0.83	0.68
MapleCaDiCaL_LBD-990_275	0.81	0.71
MapleCaDiCaL_LBD-990_500	0.84	0.72
MapleCaDiCaL_PPD-500_500	0.82	0.71
MapleCaDiCaL_PPD-950_950	0.84	0.73
MergeSat-bve_gates	0.75	0.71
MergeSat-bve_semgates	0.74	0.73
MergeSat-thread1	0.68	0.67
MiniSat+XorEngine	0.79	0.77
PreLearn-kissat-PreLearn-kissat.sh	0.66	0.54
PreLearn-kissat-PreLearn-tern-kissat.sh	0.55	0.46
ReEncode-kissat-ReEncode-pair-kissat.sh	0.75	0.68
SBVA-sbva_cadical	0.74	0.55
SBVA-sbva_kissat	0.78	0.65
SeqFROST	0.81	0.73
SeqFROST-ERE-All	0.81	0.71
SeqFROST-NoExtend	0.79	0.74
hKis-psids	0.81	0.72
hKis-sat_psids	0.80	0.79
hKis-unsat	0.86	0.69
hKissatInc-unsat	0.85	0.77
kissat-3.1.0	0.88	0.76
kissat-hywalk-exp	0.85	0.70
kissat-hywalk-exp-gb	0.89	0.74
kissat-hywalk-gb	0.88	0.72
kissat_inesp	0.95	0.80
tabularasat-1.0.0	0.88	0.76

Figure 8 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver.

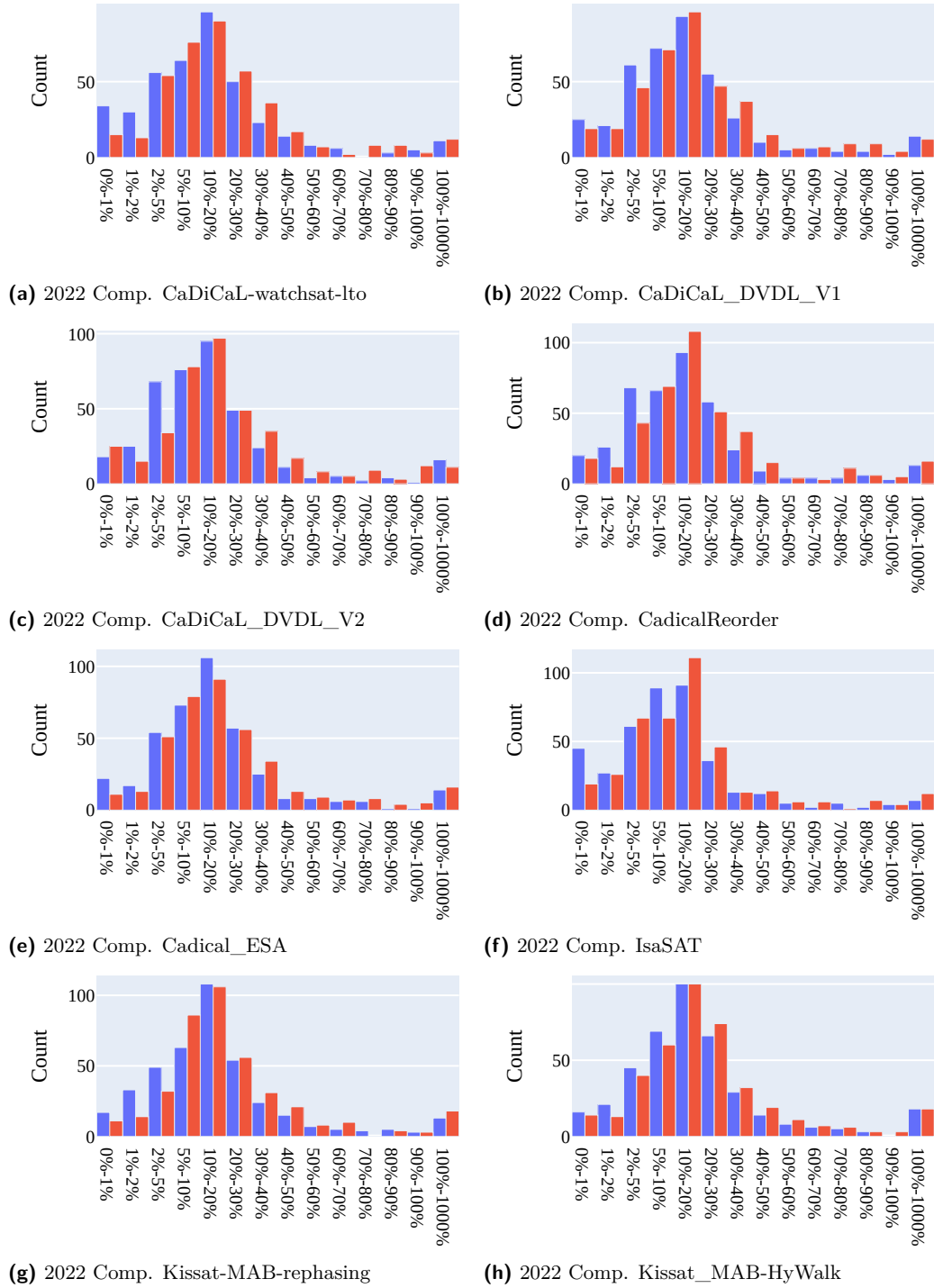


Figure 9 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

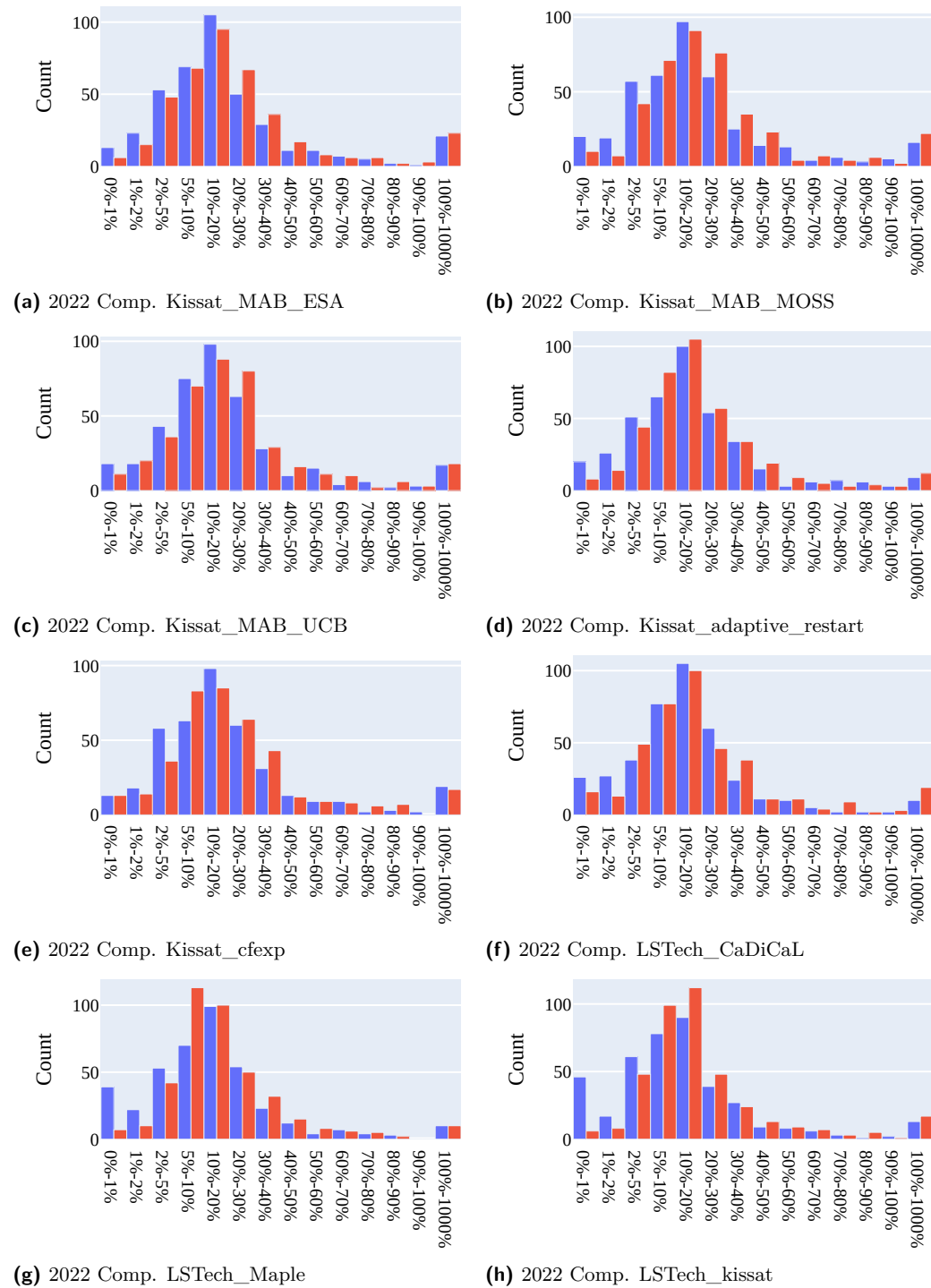


Figure 10 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATzilla 2012; in red) and the new tool (SATzilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

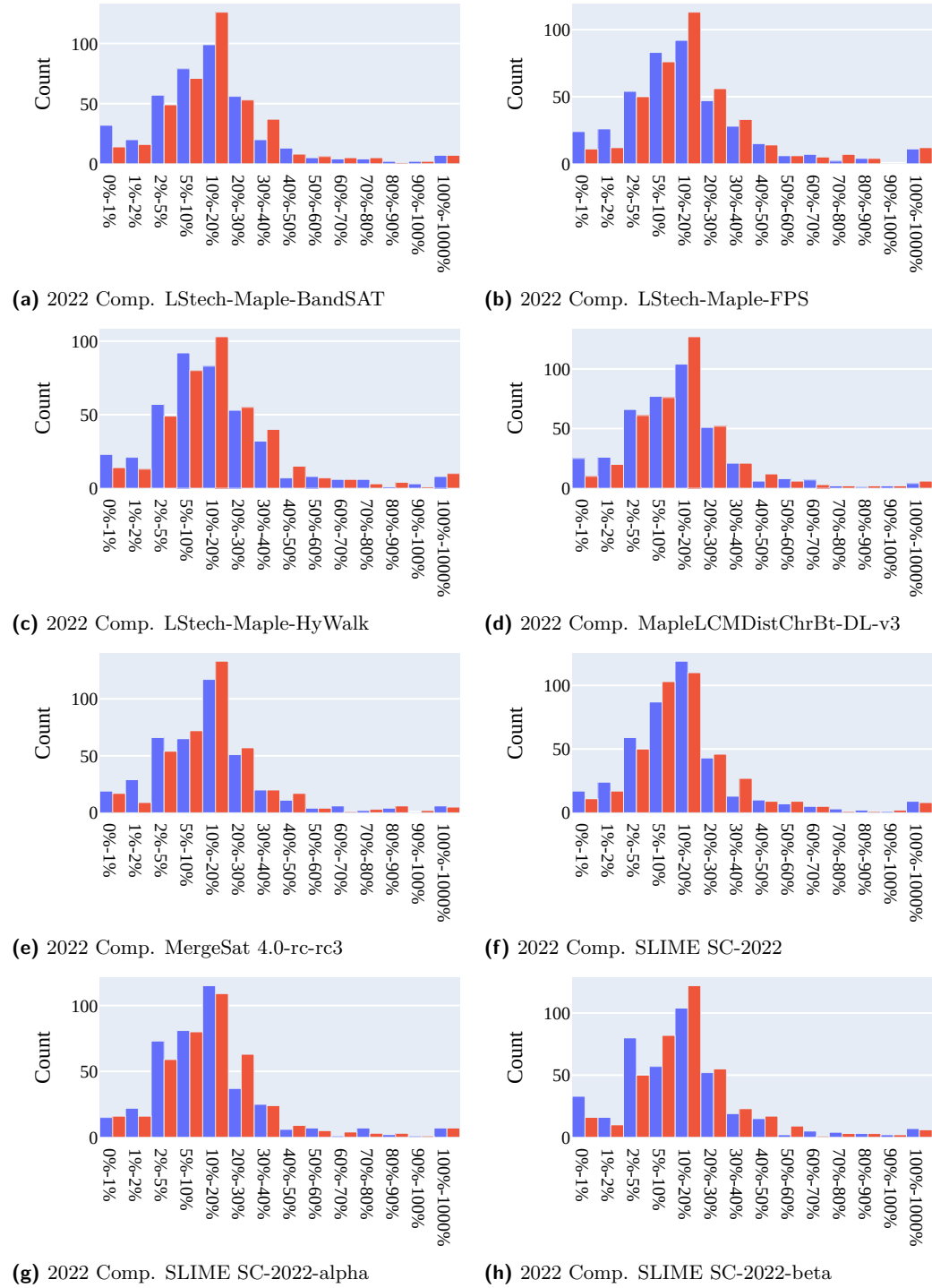


Figure 11 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

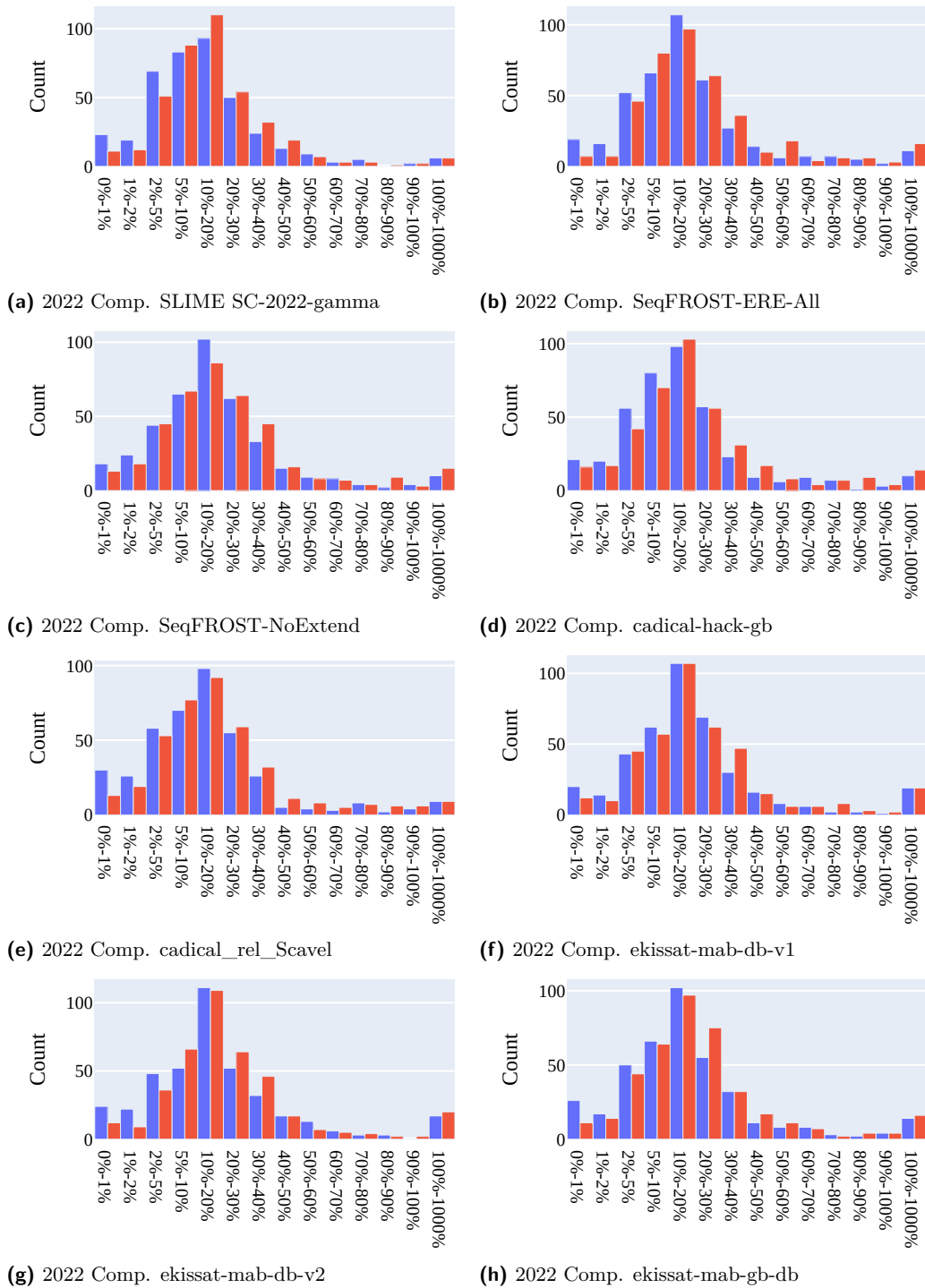


Figure 12 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

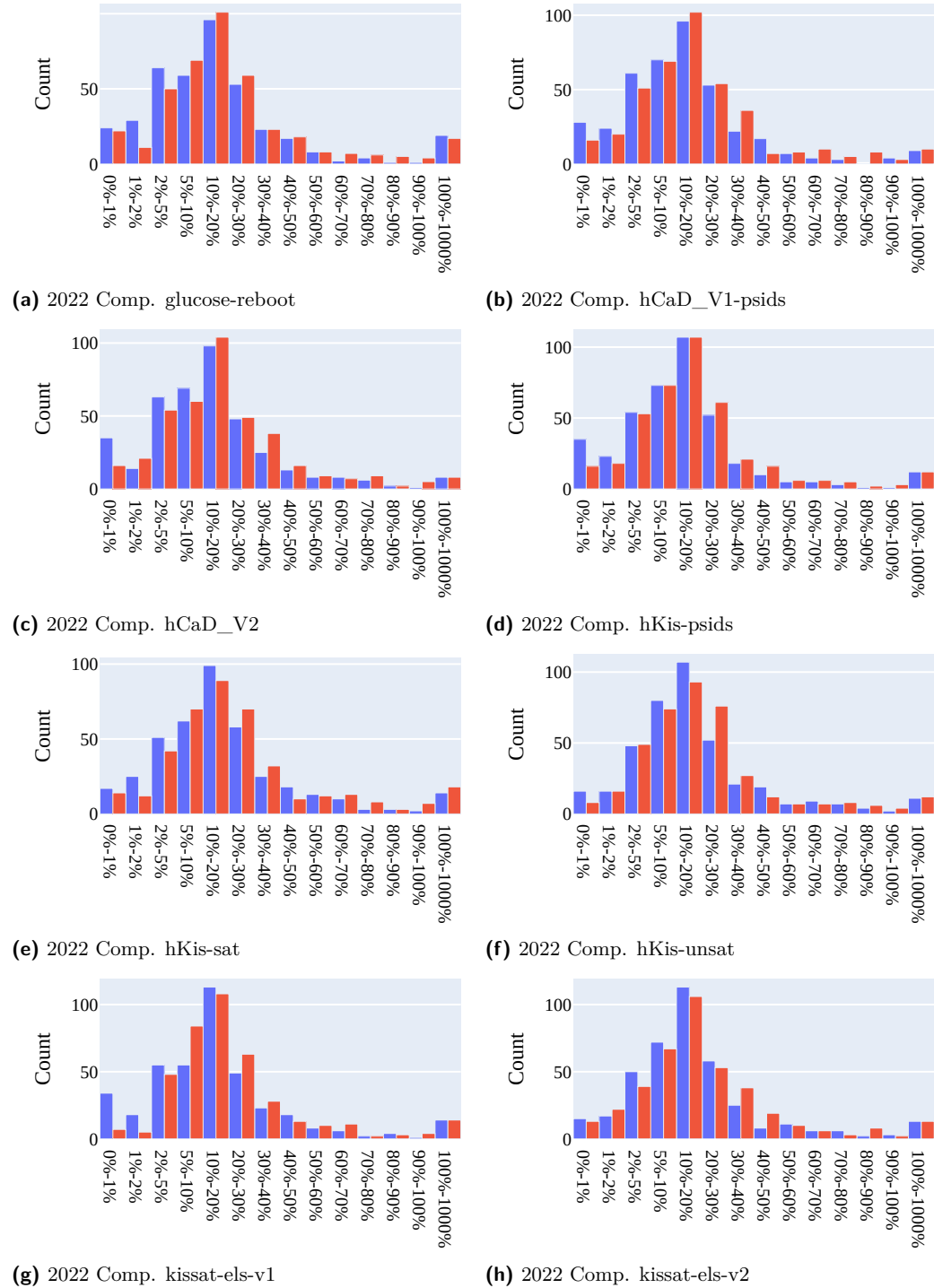


Figure 13 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

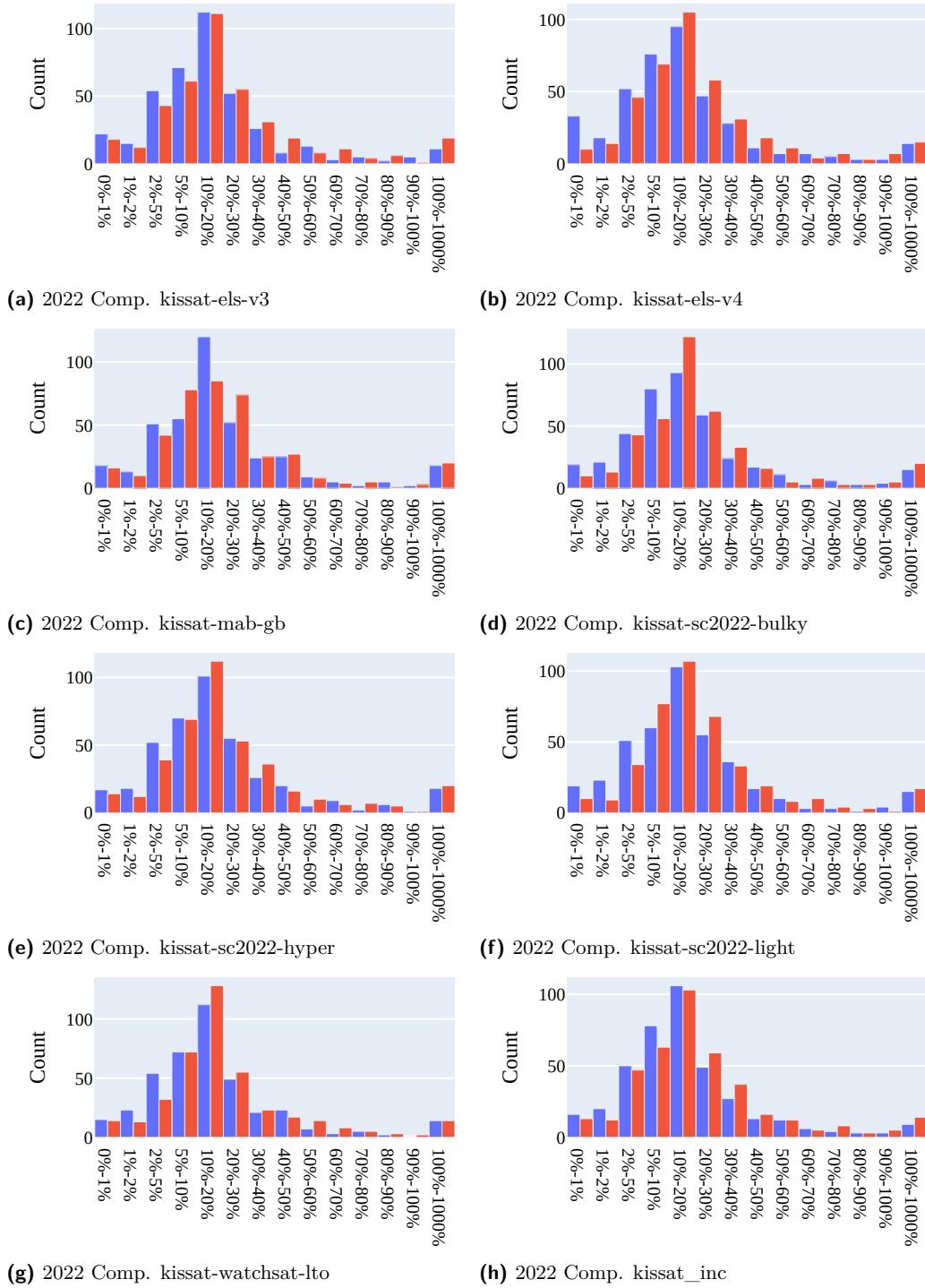


Figure 14 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

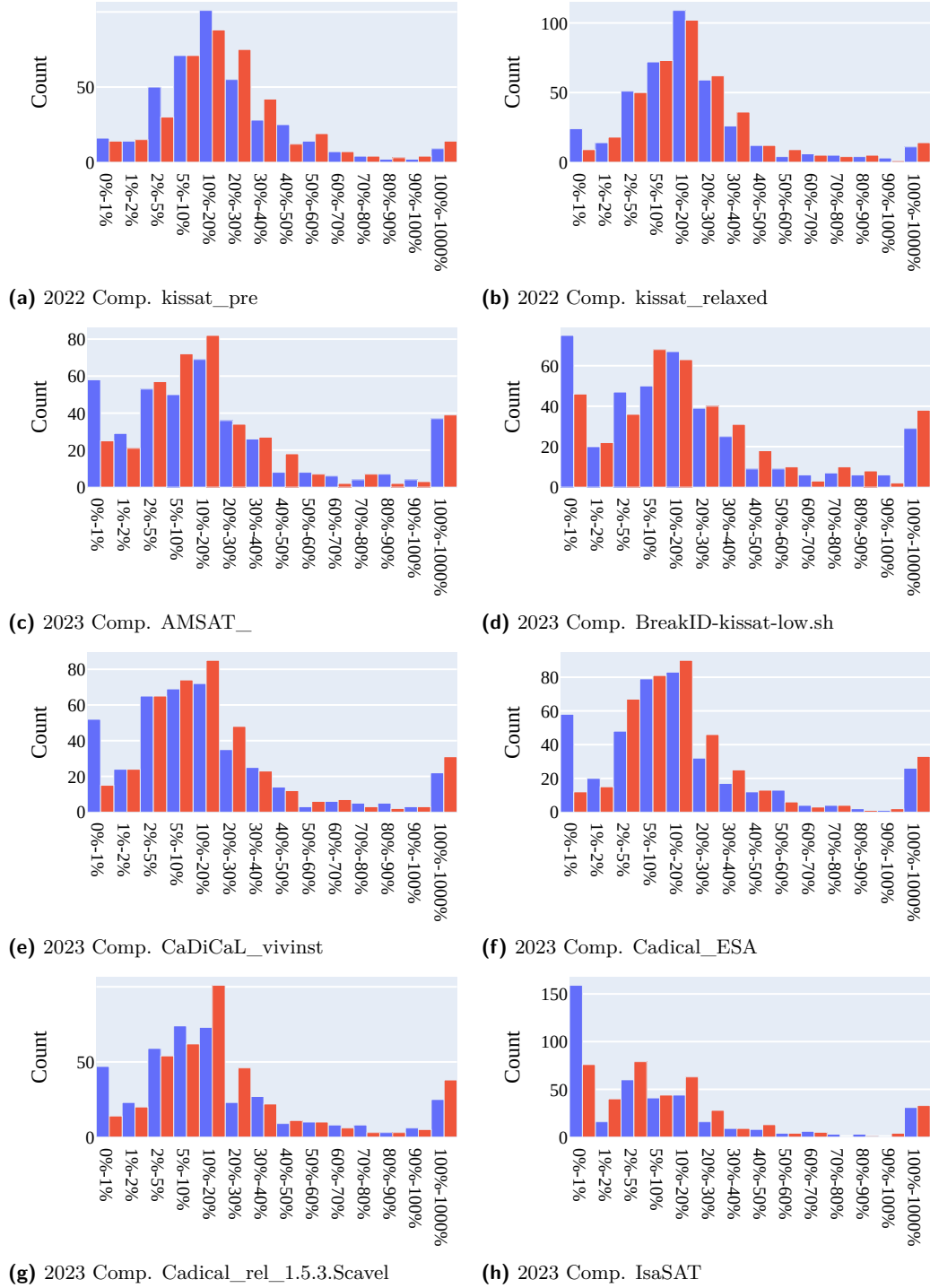


Figure 15 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

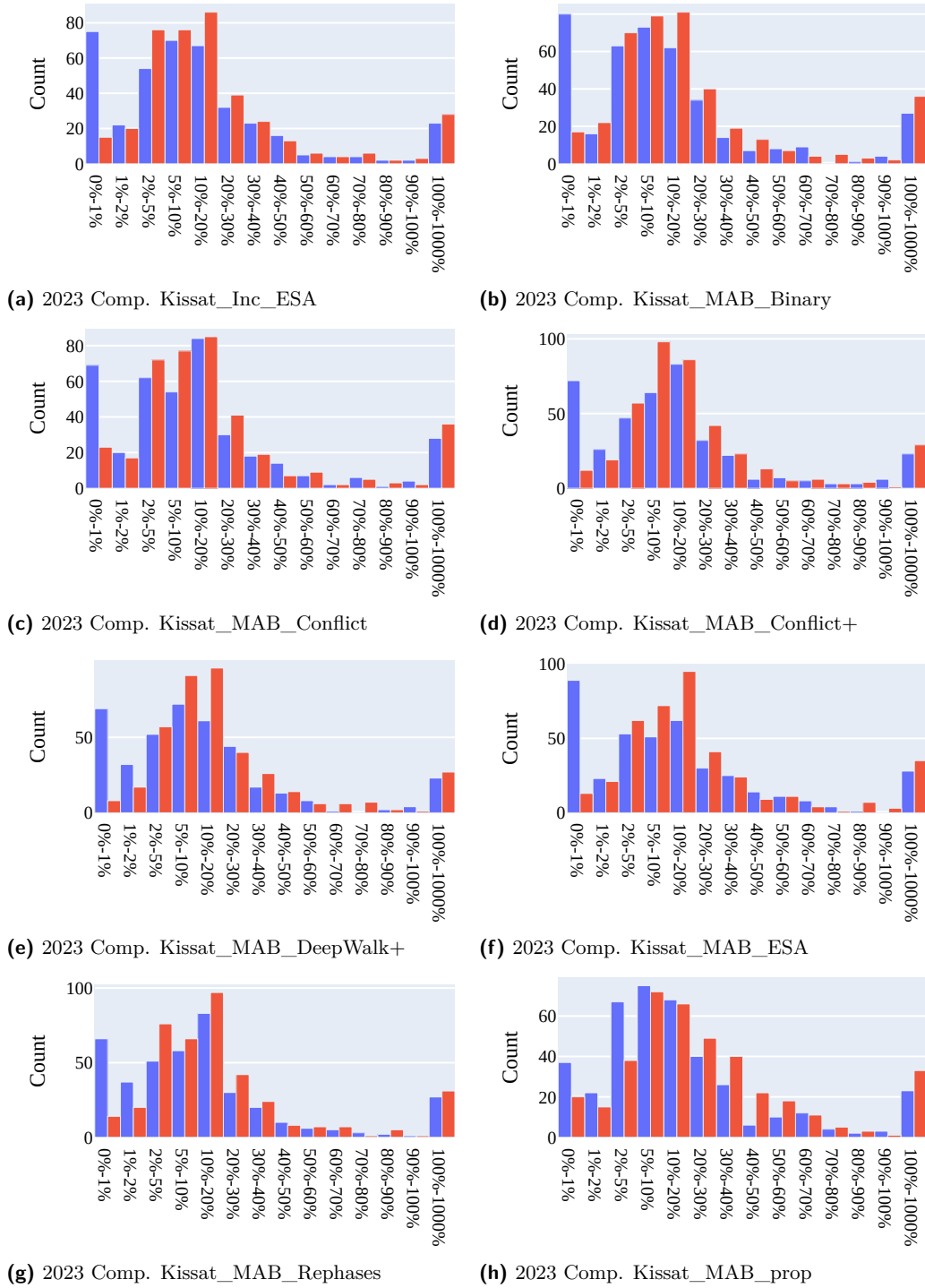


Figure 16 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATzilla 2012; in red) and the new tool (SATzilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

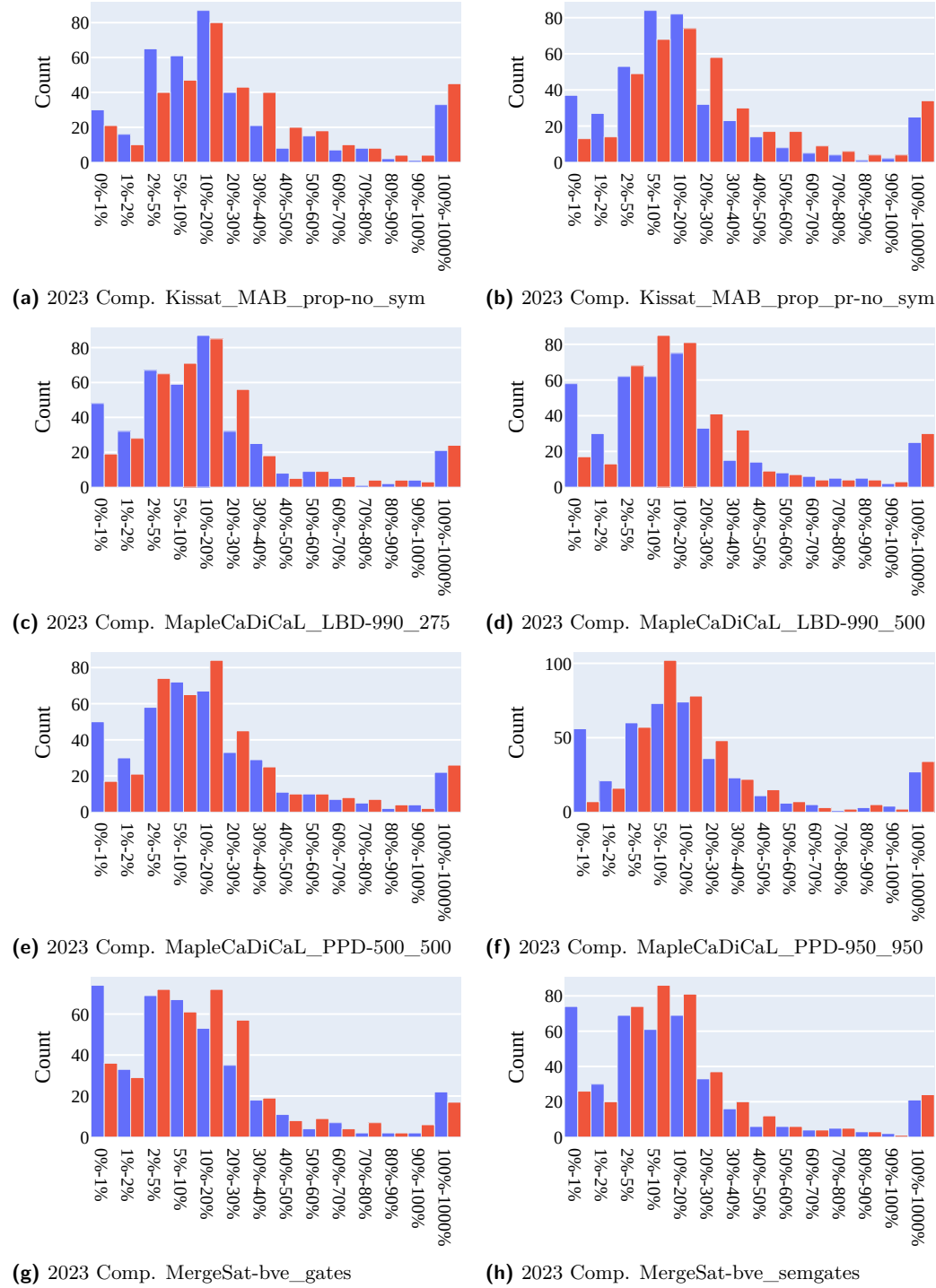


Figure 17 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

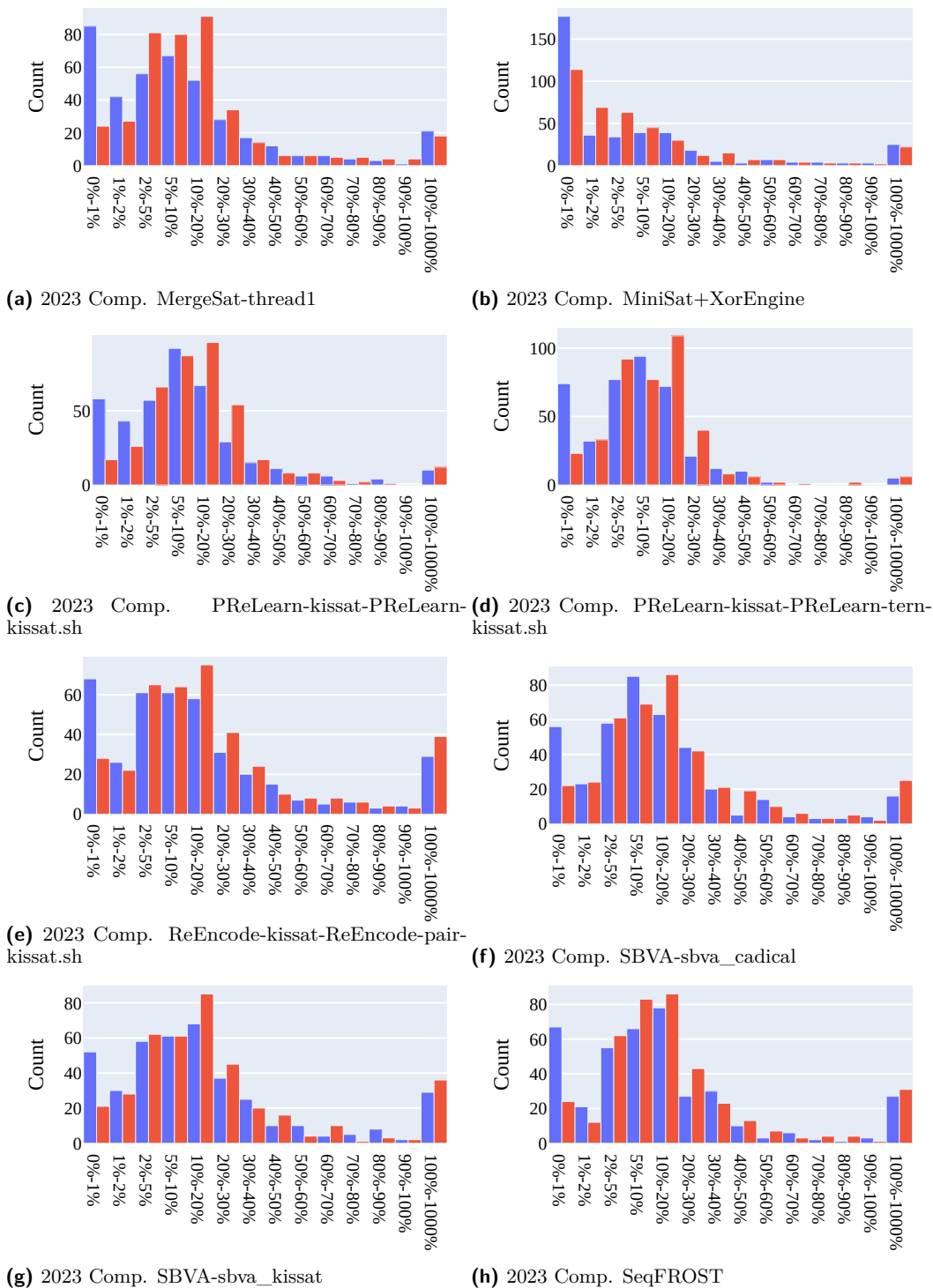


Figure 18 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATzilla 2012; in red) and the new tool (SATzilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

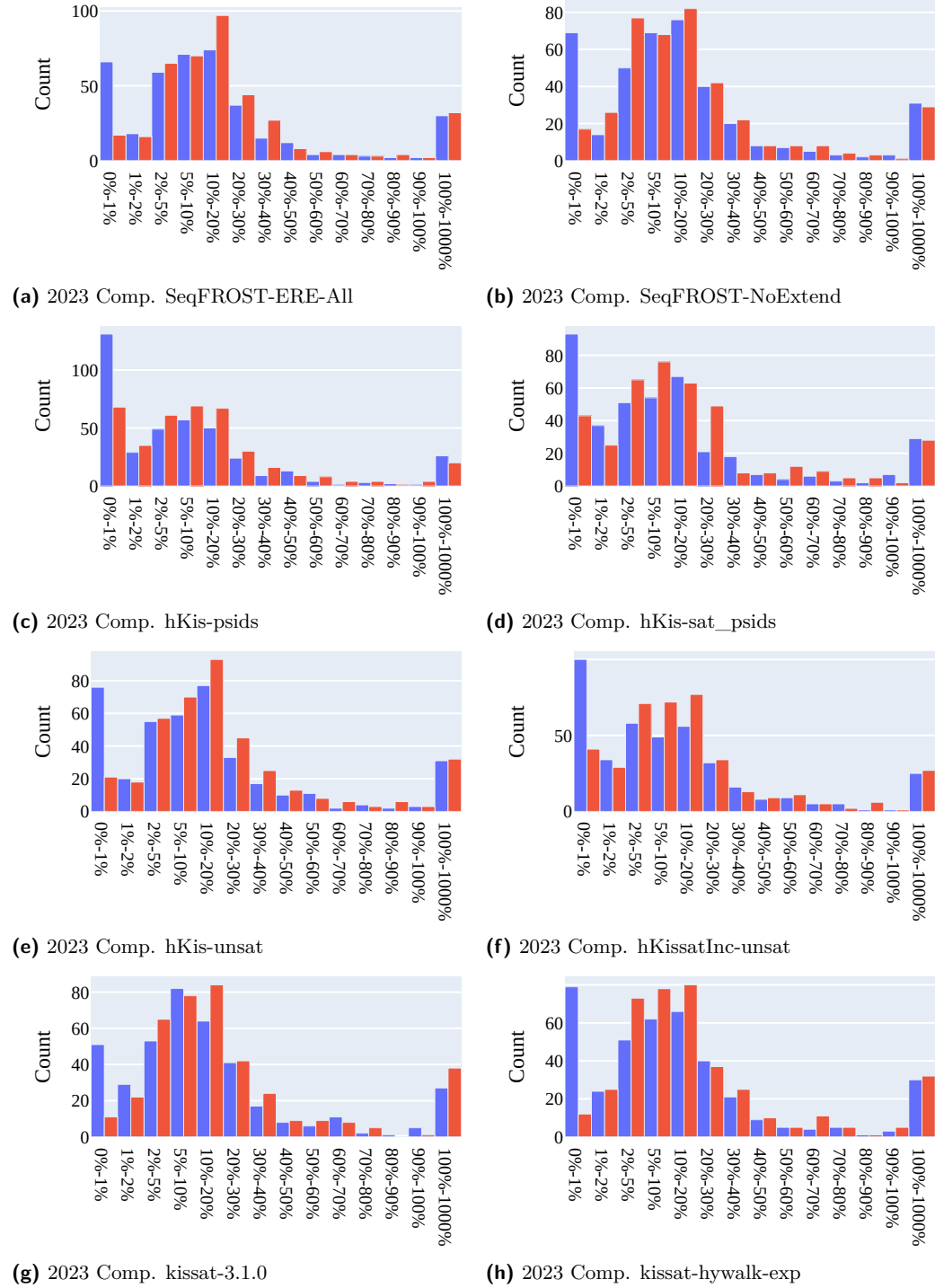


Figure 19 Histogram of the error percentage of the root mean square error (RMSE) of (log-transformed) running time prediction using a random forest with features extracted by the old (SATZilla 2012; in red) and the new tool (SATZilla 2024; in blue), on SAT solvers from the 2022 and 2023 SAT Competitions. Results are presented per solver (contd.).

